THIS NEWSLETTER DESCRIBES A FEW MINOR CHANGES TO THE SYNTAX AND SEMANTICS OF THE SETL SYSTEM NOW UNDER DEVELOPMENT.

CODE BLOCKS

WE ARE RESTORING A FORM OF #CODE BLOCKS# SIMILAR TO THAT DESCRIBED IN #ON PROGRAMMING#. A CODE BLOCK IS A SERIES OF STATEMENTS WHICH RETURNS A VALUE, AND CAN THUS BE USED WHEREVER AN EXPRESSION IS EXPECTED.

THE FORM OF A CODE BLOCK IS:

BEGIN. <STATEMENT LIST> END.;

STATEMENTS OF THE FORM

RETURN <EXPRESSION>;

CAN APPEAR OMLY WITHIN CODE BLOCKS. THEIR ACTION IS TO EXIT THE CODE BLOCK, RETURNING. <EXPRESSION> AS THE VALUE OF THE BLOCK. THIS IS THE ONLY LEGAL WAY TO EXIT A CODE BLOCK.

SUBROUTINES, FUNCTIONS, AND USER DEFINED OPERATORS

A SUBROUTINE DEFINITION CONSISTS OF A DEFINE STATEMENT, A GROUP OF DECLARATIONS, A STATEMENT LIST, AND AN END STATEMENT.

A FUNCTION OR OPERATOR DEFINITION CONSISTS OF A DEFINEF STATEMENT, A GROUP OF DECLARATIONS, AN EXPRESSION, AND AN END STATEMENT. THE EXPRESSION MAY OF COURSE BE A CODE BLOCK.

WE ALLOW NILADIC OPERATOR DEFINITIONS, FOR EXAMPLE:

DEFINEF MYTNEWAT.;
NEWAT;
END:

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ITERATORS

ANY ITERATOR MAY NOW HAVE A WHILE CONDITION; THE CURRENT WHILE ITERATOR IS ELIMINATED. THUS THE SYNTAX OF ITERATORS BECOMES:

<RANGE PART> > SET ITERATOR, MAP ITERATOR OR ARITHMETIC ITERATOR
<WHILE PART> > WHILE. <EXPRESSION>

SOME EXAMPLES ARE:

(× I = 1, 2 ... WHILE I LT. 50) PRINT I;;

 $S = \leq \langle X, Y \rangle \Rightarrow S1 \text{ WHILE } X \text{ LT. } Y \uparrow C(X, Y) \geq ;$

COMPOUND ITERATORS MAY HAVE MORE THAN ONE WHILE CONDITION, FOR EXAMPLE,

(* X = T(I) WHILE X NE. DM., Y = F(X) WHILE. Y EQ. O) PRINT X;;

HERE THE OUTER LOOP CONTINUES WHILE X IS NOT OMEGA; THE INNER LOOP
CONTINUES WHILE Y IS ZERO.

TRUE AND FALSE

AFTER MUCH DEBATE ABOUT THE MEANING OF TRUE AND FALSE, WE HAVE COME TO THE FOLLOWING CONCLUSIONS:

≠TRUE≠ IS DEFINED TO BE THE BIT STRING B≠1≠, AND ≠FALSE≠ IS DEFINED TO BE THE BIT STRING E≠0≠.

IF AN EXPRESSION IS USED AS THE CONDITION OF AN IF OR WHILE STATEMENT, ITS VALUE MUST BE IN THE SET \$3\pm 1\pm, \$b\pm 0\pm 2. ANY OTHER VALUE WILL CAUSE AN ABORT.

RANGE APPLICATION AND LEXICAL CHANGES

WE ARE CONSIDERING MAKING A SERIES OF LEXICAL CHANGES WHICH WILL MAKE THE LANGUAGE MORE READABLE. AS A SIDE EFFECT OF THESE CHANGES, SOME OF THE GENERALIZED RANGE CONSTRUCTS WILL MAVE TO BE ELIMINATED.

TUPLE FORMERS ARE NOW WRITTEN USING SQUARE BRACKETS, 1.8. [1, 2].

THE RELATIONAL OPERATORS ARE NOW WRITTEN:

OLD	NEW
LT.	<
LE.	< =
GT.	>
GE.	> z
EQ.	, ` =
NE.	/=

ASSIGNMENTS ARE WRITTEN #1=#. THERE ARE TWO ON THE FLY ASSIGNMENT OPERATORS WHICH REPLACE THE #IS# OPERATOR. #A 1= B# ASSIGNS B TO A AND HAS THE VALUE B; #B =: A# DOES LIKEWISE. THE TARGET OF AN ON THE FLY ASSIGNMENT MUST ALWAYS BE A SIMPLE NAME.

THE RANGE CONSTRUCT CAN ONLY BE USED IN THE CONTEXTS FISH AND FISH, S2, ... SNI, WHERE F IS A MAP, FUNCTION, OR SUBROUTINE.

CASE EXPRESSIONS

WE INTRODUCE A NEW CASE EXPRESSION, WHICH IS PARTICULARLY HANDY FOR USE IN BACKTRACKING. THIS EXPRESSION HAS THE FORM

N CASE. (<LIST OF EXPRESSIONS>)

ITS ACTION IS TO EVALUATE THE N-TH EXPRESSION IN THE LIST OF EXPRESSIONS AND RETURN ITS VALUE. THE REMAINING EXPRESSIONS ARE NOT EVALUATED.

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O BACKTRACKING CONSTRUCTS

THIS SECTION GIVES THE SYNTAX OF THE VARIOUS BACKTRACKING PRIMITIVES DISCUSSED IN NL. 186.

THE BACKTRACKING DECLARATIONS ARE:

BACK. <NAMELIST> END; NOBACK. <NAMELIST> END.;

WE ADD THE NILADIC OPERATORS OK AND SUCCESS, AND THE STATEMENTS FAIL, ACCEPT AND REJECT.

THE NONDETERMINISTIC ARB OPERATION IS REPRESENTED BY A BACKWARDS EPSILON, WRITTEN *% *. THE DETERMINISTIC ARB IS WRITTEN *ARB*.

NONDETERMINISTIC CASE EXPRESSIONS MAY BE WRITTEN

WN CASE. (EXPl. EXP2, ... EXPN)

OR ABBREVIATED

% CASE. (EXPl, EXP2, ... EXPN)